RESEARCH
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ABSTRACT

Search algorithms are often categorized by their node expansion strategy. One option is the depth-first strategy, a simple backtracking strategy that traverses the search space in the order in which successor nodes are generated. An alternative is the best-first strategy, which was designed to make it possible to use domain-specific heuristic information. By exploring promising parts of the search space first, best-first algorithms are usually more efficient than depth-first algorithms.

In programs that play minimax games such as chess and checkers, the efficiency of the search is of crucial importance. Given the success of best-first algorithms in other domains, one would expect them to be used for minimax games too. However, all high-performance game-playing programs are based on a depth-first algorithm.

This study takes a closer look at a depth-first algorithm, Alpha-Beta, and a best-first algorithm, SSS*. The prevailing opinion on these algorithms is that SSS* offers the potential for a more efficient search, but that its complicated formulation and exponential memory requirements render it impractical. The theoretical part of this work shows that there is a surprisingly straightforward link between the two algorithms—for all practical purposes, SSS* is a special case of Alpha-Beta. Subsequent empirical evidence proves the prevailing opinion on SSS* to be wrong: it is not a complicated algorithm, it does not need too much memory, and it is also *not* more efficient than depth-first search.

Over the years, research on Alpha-Beta has yielded many enhancements, such as transposition tables and minimal windows with re-searches, that are responsible for the success of depth-first minimax search. The enhancements have made it possible to use a depth-first procedure to expand nodes in a best-first sequence. Based on these insights, a new algorithm is presented, MTD(f), which out-performs both SSS* and NegaScout, the Alpha-Beta variant of choice by practitioners.

In addition to best-first search, other ways for improvement of minimax search algorithms are investigated. The tree searched in Alpha-Beta's best case is usually considered to be equal to the minimal tree that has to be searched by any algorithm in order to find and prove the minimax value. We show that in practice this assumption is not valid. For non-uniform trees, the real minimal tree—or rather, graph—that proves the minimax value is shown to be significantly smaller than Alpha-Beta's best case. Thus, there is more room for improvement of full-width minimax search than is generally assumed.

Preface

I would like to thank my advisors Arie de Bruin, Wim Pijls, and Jonathan Schaeffer, for being just as eager as I to understand minimax trees. We have spent countless hours discussing trees and algorithms, trying to find the hidden essence. Their experience with minimax search, as well as their different backgrounds, were of great value. They know how much this research owes to their willingness to spend so much time discussing ideas, combining theory and practice, trying to understand trees. It has been a great time, for which I thank you three.

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From October 1994 to October 1995 I spent most of my time at the University of Alberta in Edmonton, Canada, with Jonathan Schaeffer. Working together with someone with so much experience and energy, and so many ideas, has been an extraordinary experience. Jonathan, I thank you for giving me the opportunity to do research with you and to learn from you; but most of all, I thank you for good times. In Edmonton many people did their best to make my stay as pleasant as possible. Jonathan and Stephanie Schaeffer made sure that I felt at home, helping me wherever and whenever they could. I thank the people of the department of Computing Science of the University of Alberta for creating a good working environment. I have had many discussions on minimax search with Mark Brockington, author of Keyano, one of the top Othello programs. Yngvi Björnsson and Andreas Junghanns, who started building a new chess program, were always keen to discuss new ideas too. Rob Lake answered my troff questions. I thank all of the members of the Games research group for making both the scheduled and the informal meetings interesting and enjoyable.

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