

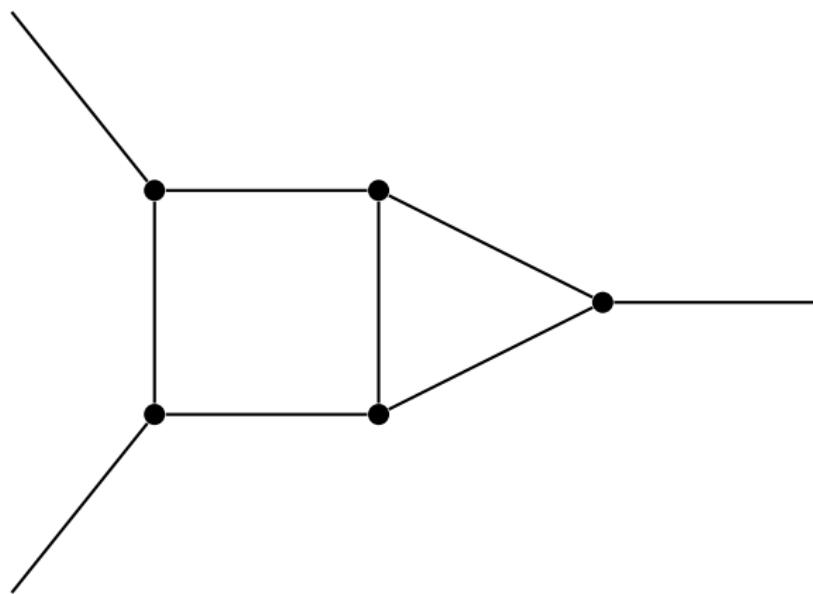
Tokové polynómy k -pólov

Diplomová práca

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doc. RNDr. Robert Lukočka, PhD.

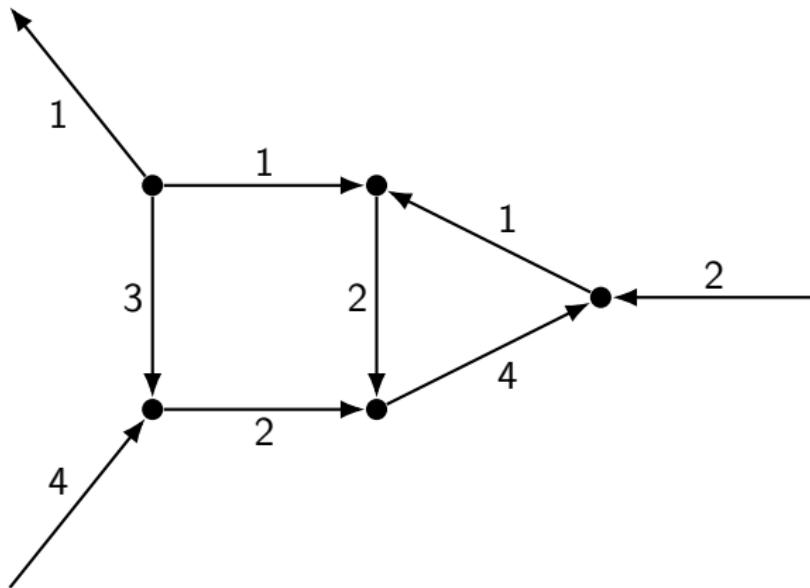
Katedra Informatiky
Fakulta Matematiky, Fyziky a Informatiky
Univerzita Komenského

k -pól



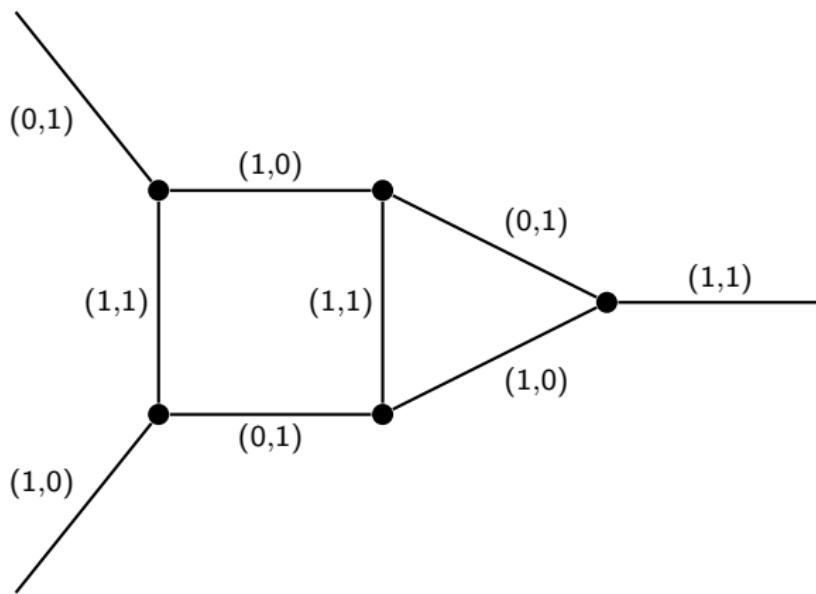
Obr.: Príklad kubického 3-pólu.

Nikde-nulový tok nad grupou



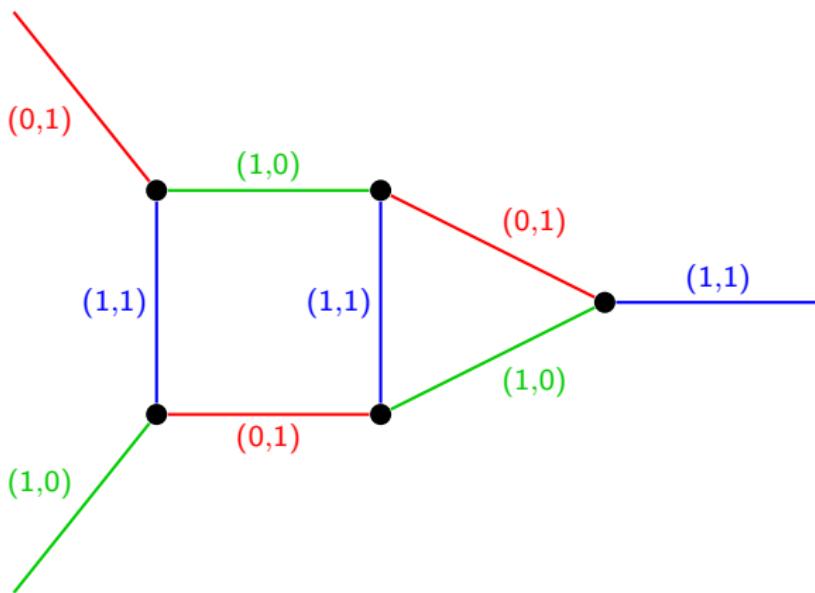
Obr.: Príklad nikde-nulového \mathbb{Z}_5 -toku v 3-póle.

$(\mathbb{Z}_2 \times \mathbb{Z}_2)$ -tok



Obr.: Príklad nikde-nulového $(\mathbb{Z}_2 \times \mathbb{Z}_2)$ -toku v 3-póle. Na orientácii hrán nezáleží – v $(\mathbb{Z}_2 \times \mathbb{Z}_2, +)$ je každý prvok involúciou.

$(\mathbb{Z}_2 \times \mathbb{Z}_2)$ -tok



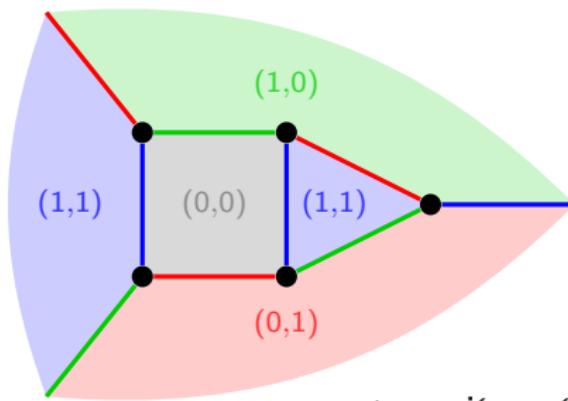
Obr.: Nikde-nulový $(\mathbb{Z}_2 \times \mathbb{Z}_2)$ -tok v kubickom k -póle tvorí hranové 3-farbenie.

Tokovo-farbiaca dualita

Planárny k -pól G :

m -farbenie oblastí \Leftrightarrow nikde-nulový m -tok

veta o 4 farbách \Leftrightarrow nikde-nulový $(\mathbb{Z}_2 \times \mathbb{Z}_2)$ -tok

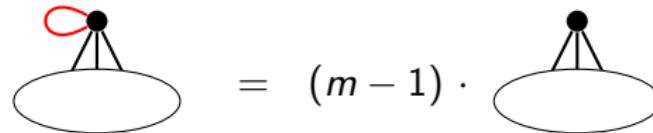


+ zaujímavé sú kubické grafy

Počet tokov $f(G)$

- ak e je slučka v G :

$$f(G) = (m - 1) \cdot f(G - e)$$



- inak:

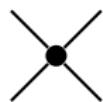
$$f(G) = f(G / e) - f(G - e)$$



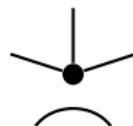
Základné k -póly = bez vnútorných hrán



Obr.: Planárny 3-pól bez vnútorných hrán.



Obr.: Planárne 4-póly bez vnútorných hrán.



Obr.: Planárne 5-póly bez vnútorných hrán.

$f(G)$ ako kombinácia základných k -pólov

$$\begin{array}{c} \text{Diagram of a } k\text{-pole with 4 vertices and 6 edges} \\ = a \cdot \text{Diagram of a } k\text{-pole with 1 vertex and 3 edges} + b_1 \cdot \text{Diagram of a } k\text{-pole with 2 vertices and 3 edges} \\ + b_2 \cdot \text{Diagram of a } k\text{-pole with 0 vertices and 2 edges} = \\ = 0 \cdot \text{Diagram of a } k\text{-pole with 1 vertex and 3 edges} + 1 \cdot \text{Diagram of a } k\text{-pole with 0 vertices and 2 edges} \end{array}$$

$f(G)$ ako kombinácia základných k -pólov

$$\begin{array}{c} \text{Diagram of a } k\text{-pole} \\ \text{with 4 vertices and 6 edges} \end{array} = a \cdot \begin{array}{c} \text{Diagram of a } k\text{-pole} \\ \text{with 3 vertices and 4 edges} \end{array} + b_1 \cdot \left(\begin{array}{c} \text{Diagram of a } k\text{-pole} \\ \text{with 2 vertices and 2 edges} \end{array} \right) + b_2 \cdot \begin{array}{c} \text{Diagram of a } k\text{-pole} \\ \text{with 1 vertex and 1 edge} \end{array} = \\ = 0 \cdot 1 + 1 \cdot 1 + 1 \cdot 0 \end{array}$$

$f(G)$ ako kombinácia základných k -pólov

$$\begin{array}{c} \text{Diagram of a graph } G \\ \text{with 5 vertices and 6 edges.} \\ = a \cdot \text{Diagram of a } k\text{-pole} \\ + b_1 \cdot \text{Diagram of a } k\text{-pole} \\ + b_2 \cdot \text{Diagram of a } k\text{-pole} \\ = \\ = 0 \cdot 1 + 1 \cdot 1 + 1 \cdot 0 \end{array}$$

Na \mathbb{Z}_5 :

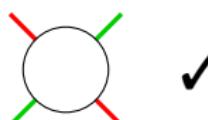
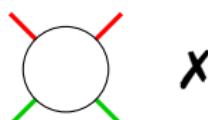
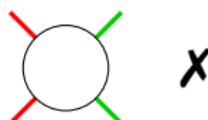
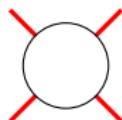
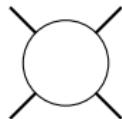
Kochol (2004). Reduction of the 5-Flow Conjecture to cyclically 6-edge-connected snarks.

Kochol (2005). Decomposition formulas for the flow polynomial.

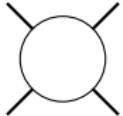
Kochol (2006). Restrictions On Smallest Counterexamples To The 5-Flow Conjecture.

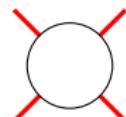
Kochol (2010). Smallest counterexample to the 5-flow conjecture has girth at least eleven.

Motivácia

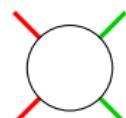


Motivácia


$$= a \cdot \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} + b_1 \cdot \left(\begin{array}{c} \diagup \\ \diagdown \end{array} \right) + b_2 \cdot \begin{array}{c} \diagdown \\ \diagup \end{array}$$

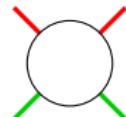


$$a + b_1 + b_2 \geq 0$$



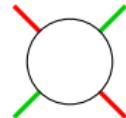
X

$$a + b_1 = 0$$



X

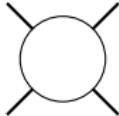
$$a + b_2 = 0$$

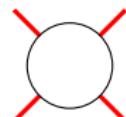


✓

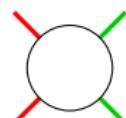
$$a > 0$$

Motivácia


$$= a \cdot \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} + b_1 \cdot \left(\begin{array}{c} \diagup \\ \diagdown \end{array} \right) + b_2 \cdot \begin{array}{c} \diagup \\ \diagdown \\ \diagup \end{array}$$

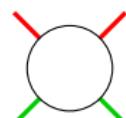


$$a + b_1 + b_2 \geq 0$$



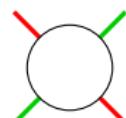
X

$$a + b_1 = 0 \quad \Rightarrow b_1 = -a$$



X

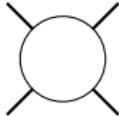
$$a + b_2 = 0 \quad \Rightarrow b_2 = -a$$

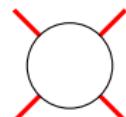


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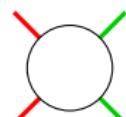
$$a > 0 \quad \Rightarrow a + b_1 + b_2 < 0 \quad \text{SPOR}$$

Motivácia


$$= a \cdot \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} + b_1 \cdot \left(\begin{array}{c} \diagup \\ \diagdown \end{array} \right) + b_2 \cdot \begin{array}{c} \diagup \\ \diagdown \\ \diagup \end{array}$$

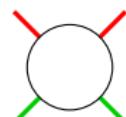


$$a + b_1 + b_2 \geq 0$$



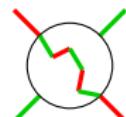
X

$$a + b_1 = 0 \quad \Rightarrow b_1 = -a$$



X

$$a + b_2 = 0 \quad \Rightarrow b_2 = -a$$



✓

$$a > 0 \quad \Rightarrow a + b_1 + b_2 < 0 \quad \text{SPOR}$$

Naša práca

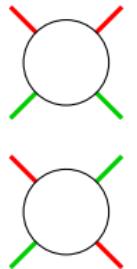
Koeficienty a farbenia

k	koeficienty	farbenia
2	1	1
3	1	1
4	3	4
5	6	10
6	15	31
7	36	91
8	91	274
...	$\Theta(3^k/k^{3/2})$	$\Theta(3^k)$

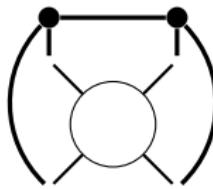
dôkaz ostrej nerovnosti pre $k \geq 4$

→ pre planárne grafy sú
koeficienty efektívnym
nástrojom na štúdium farbení
(pre neplanárne však nie)

Indikátor (p, r) -súvislosti

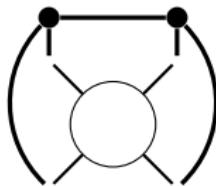
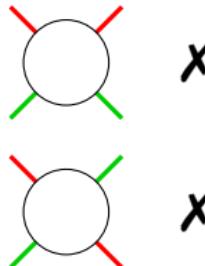


$\times \quad \times$



$$2a + b_2 = 0$$

Indikátor (p, r) -súvislosti

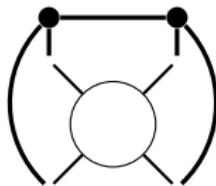
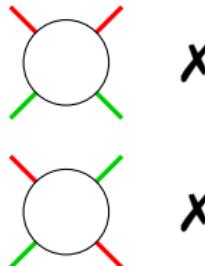


$$2a + b_2 = 0$$

\Updownarrow

G je $(2, 2)$ -nesúvislý

Indikátor (p, r) -súvislosti



$$2a + b_2 = 0$$

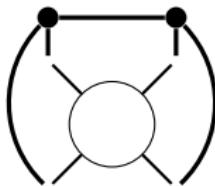
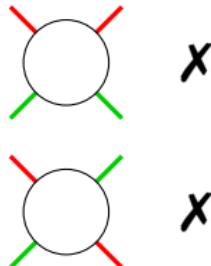
 \Updownarrow

G je $(2, 2)$ -nesúvislý

Veta

G je (p, r) -súvislý práve vtedy, keď súčet počtov (p, r) -nevyvážených farbení je nenulový.

Indikátor (p, r) -súvislosti



$$2a + b_2 = 0$$

\Updownarrow

G je $(2, 2)$ -nesúvislý

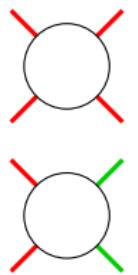
Veta

G je (p, r) -súvislý práve vtedy, keď súčet počtov (p, r) -nevyvážených farbení je nenulový.

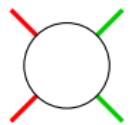
Veta

G je (p, r) -súvislý práve vtedy, keď aspoň jeden jeho (p, r) -súvislý koeficient je nenulový.

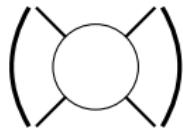
Indikátor (p, r) -2-súvislosti



\times

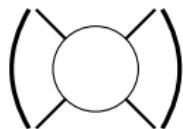
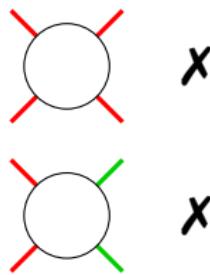


\times



$$2a + 2b_1 + b_2 = 0$$

Indikátor (p, r) -2-súvislosti



$$2a + 2b_1 + b_2 = 0 \Leftrightarrow G \text{ má } (2, 2)\text{-most}$$

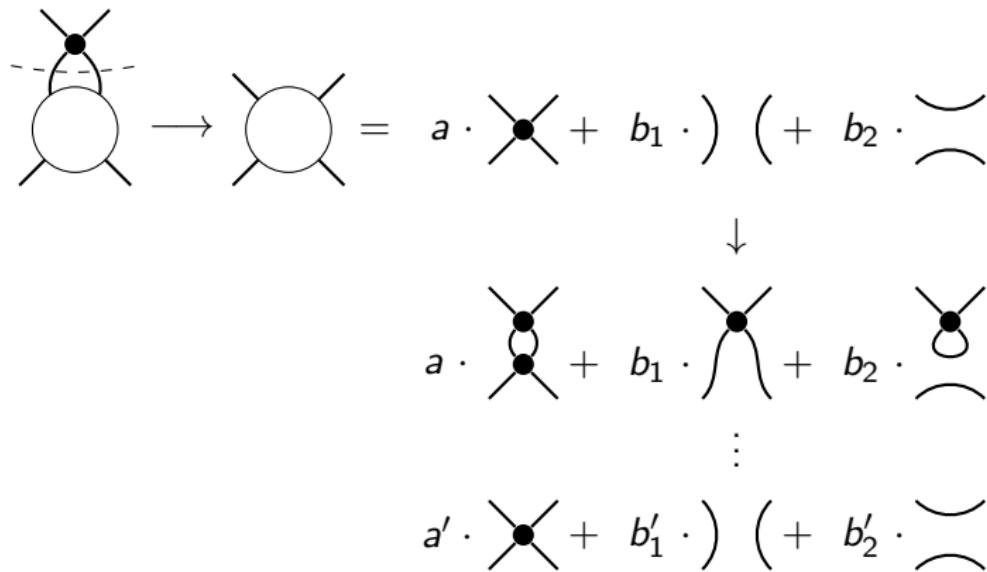
Veta

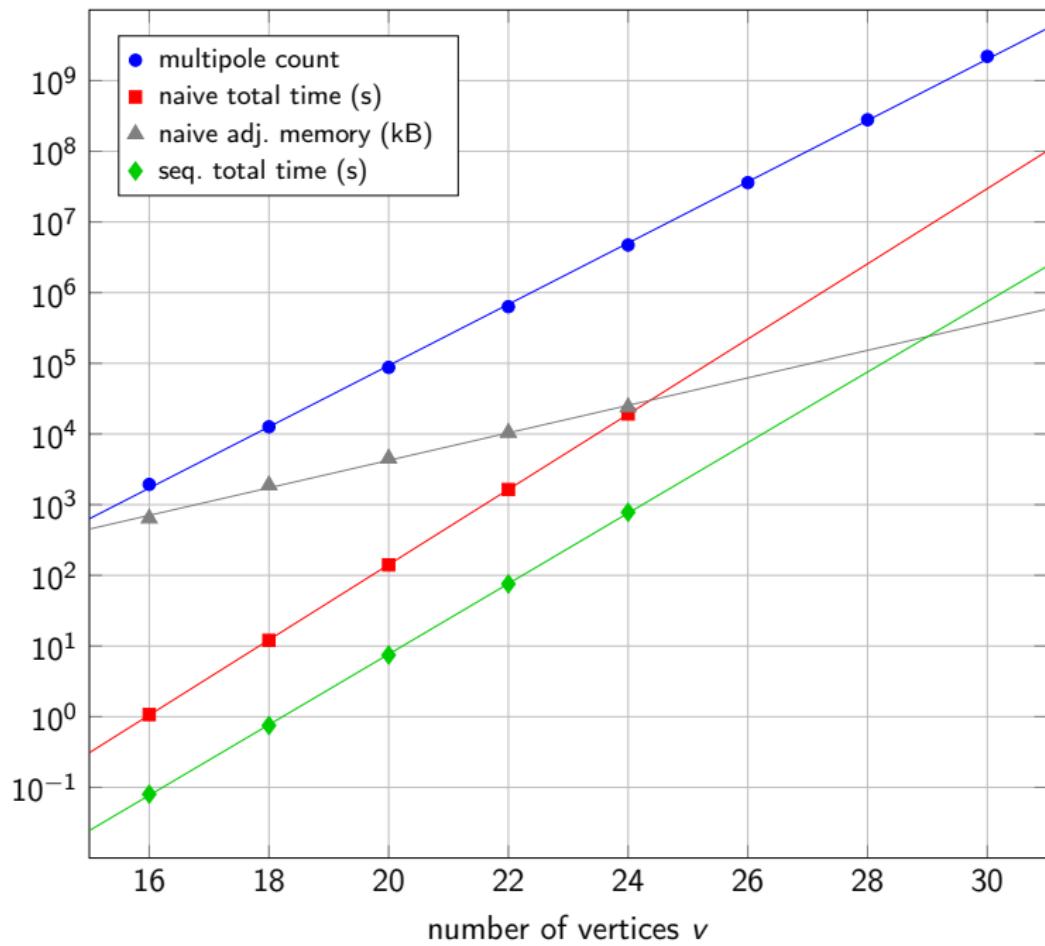
G je (p, r) -2-súvislý práve vtedy, keď súčet počtov (p, r) -vyvážených farbení je nenulový.

Algoritmus výpočtu koeficientov

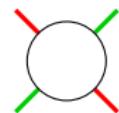
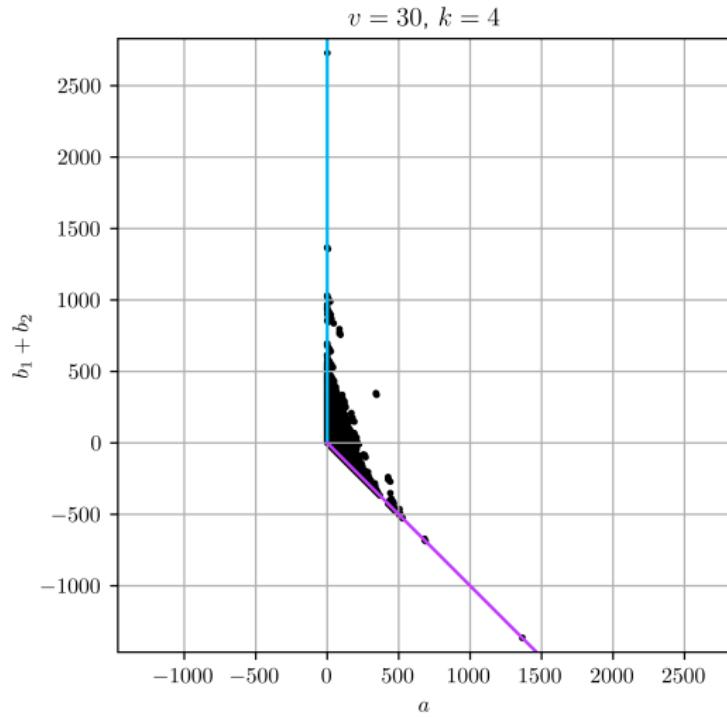
- NAIVE: rekurzia, memoizácia, orezávanie

- SEQ:

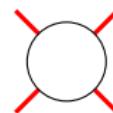




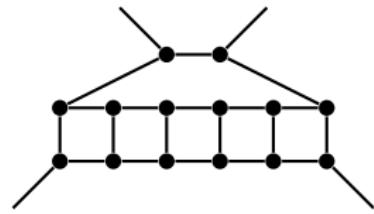
Teoretické a empirické obmedzenia ($k = 4$)



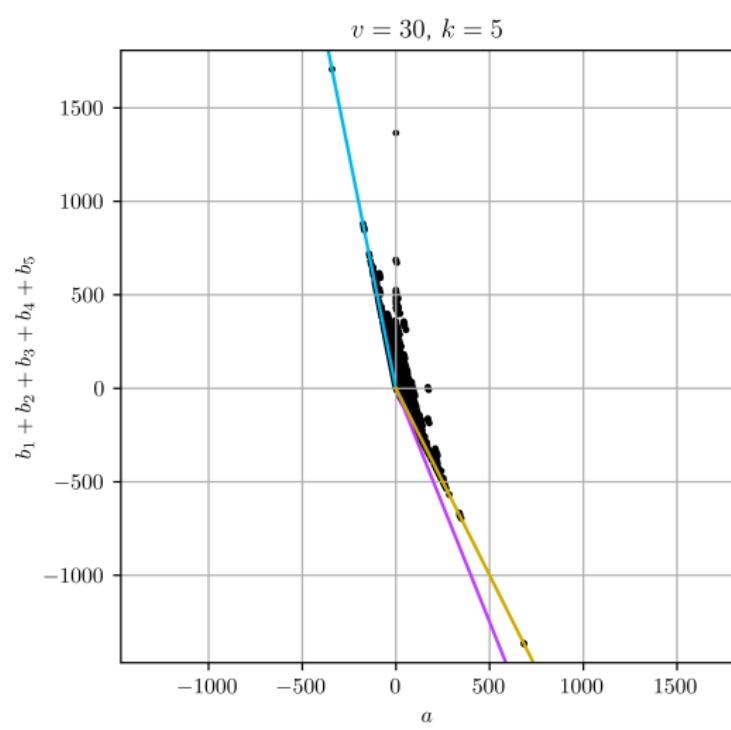
$$a \geq 0$$



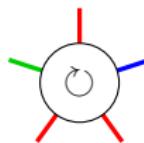
$$a + b_1 + b_2 \geq 0$$



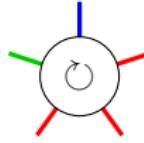
Teoretické a empirické obmedzenia ($k = 5$)



$$b = b_1 + b_2 + b_3 + b_4 + b_5$$

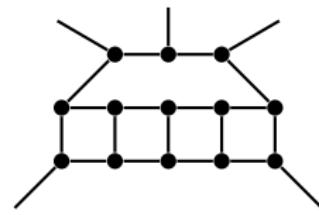


$$5a + b \geq 0$$



$$5a + 2b \geq 0$$

$$2a + b \stackrel{?}{\geq} 0$$



Teoretické a empirické obmedzenia ($k = 5$)

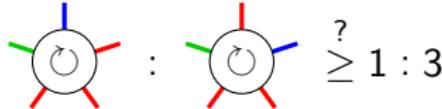
$$5a + 2b = \frac{1}{3}(5a + b) + \underbrace{\frac{5}{3}(2a + b)}_{\geq 0}$$

$$\frac{5a + 2b}{5a + b} \stackrel{?}{\geq} \frac{1}{3}$$

Teoretické a empirické obmedzenia ($k = 5$)

$$5a + 2b = \frac{1}{3}(5a + b) + \underbrace{\frac{5}{3}(2a + b)}_{\geq 0}$$

$$\frac{5a + 2b}{5a + b} \stackrel{?}{\geq} \frac{1}{3}$$


$$\stackrel{?}{\geq} 1 : 3$$

Hypotéza

Pre každý ofarbiteľný planárny kubický 5-pól tvoria farbenia typu 00012 aspoň $1/4$ celkového počtu farbení.

Hypotéza

Pre každý ofarbitel'ný planárny kubický 5-pól tvoria farbenia typu 00012 aspoň $1/4$ celkového počtu farbení.

Veta

Najmenší protipríklad (ak existuje):

- je súvislý,
- neobsahuje most,
- neobsahuje 2-rez *okrem dvojíc hrán susedných s trčiacou*,
- má aspoň 36 vrcholov.

$$\begin{array}{c} \text{Diagram with two vertices } G' \text{ and } G'' \\ \text{Diagram with one vertex } G'' \text{ and a loop} \\ \text{Diagram with one vertex } G'' \text{ and three outgoing lines} \\ \text{Diagram with one vertex } G'' \text{ and two loops} \end{array}
 = a' \cdot
 \begin{array}{c} \text{Diagram with one vertex } G'' \text{ and a loop} \end{array}
 + b'_1 \cdot
 \begin{array}{c} \text{Diagram with one vertex } G'' \text{ and three outgoing lines} \end{array}
 + b'_2 \cdot
 \begin{array}{c} \text{Diagram with one vertex } G'' \text{ and two loops} \end{array}$$

$$\begin{array}{c} \text{Diagram with two nodes } G' \text{ and } G'' \\ \text{Diagram with one node } G'' \end{array} = a' \cdot
 \begin{array}{c} \text{Diagram with one node } G'' \text{ and a black dot} \end{array} + b'_1 \cdot
 \begin{array}{c} \text{Diagram with one node } G'' \text{ and three lines} \end{array} + b'_2 \cdot
 \begin{array}{c} \text{Diagram with one node } G'' \text{ and two loops} \end{array}$$

$$a'a'': \quad \text{Diagram with one node and two lines} = 2 \cdot \star + \star$$

$$a'b''_1: \quad \text{Diagram with one node and two lines} = \star - \star$$

$$a'b''_2: \quad \text{Diagram with one node and two lines} = 3 \cdot \star$$

$$a'b''_3: \quad \text{Diagram with one node and two lines} = \star - \star$$

$$a'b''_4: \quad \text{Diagram with one node and two lines} = 2 \cdot \star$$

$$a'b''_5: \quad \text{Diagram with one node and two lines} = 2 \cdot \star$$

$$b'_1 a'': \quad \text{Diagram with one node and two lines} = \star$$

$$b'_1 b''_1: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_1 b''_2: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_1 b''_3: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_1 b''_4: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_1 b''_5: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_2 a'': \quad \text{Diagram with one node and two lines} = 3 \cdot \star$$

$$b'_2 b''_1: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_2 b''_2: \quad \text{Diagram with one node and two lines} = 3 \cdot \star$$

$$b'_2 b''_3: \quad \text{Diagram with one node and two lines} = \star$$

$$b'_2 b''_4: \quad \text{Diagram with one node and two lines} = 0$$

$$b'_2 b''_5: \quad \text{Diagram with one node and two lines} = 0$$

$$2a + b_1 + b_2 + b_3 + b_4 + b_5$$

$$2a + b_1 + b_2 + b_3 + b_4 + b_5$$

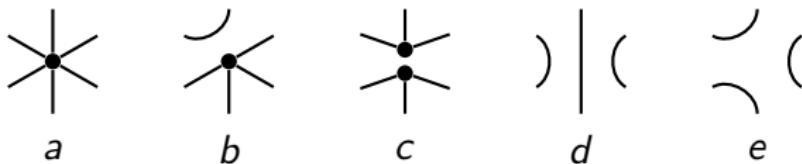
=

$$\begin{aligned} & a' (2(a'' + b_2'') + (a'' + b_4'' + b_5'')) \\ & + (a' + b_1') (2a'' + b_1'' + b_2'' + b_3'' + b_4'' + b_5'') \\ & + b_2' ((a'' + b_1'' + b_2'') + (a'' + b_2'' + b_3'') + (a'' + b_2'')) \end{aligned}$$

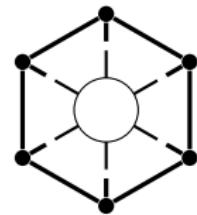
=

$$\begin{aligned} & (a' + b_1' + b_2')(2a'' + b_1'' + b_2'' + b_3'' + b_4'' + b_5'') \\ & + 2(a' + b_2')(a'' + b_2'') \\ & + (a' - b_2')(a'' + b_4'' + b_5'') \end{aligned}$$

Teoretické a empirické obmedzenia ($k = 6$)



$\circlearrowleft \beta$	a	b	c	d	e	constraint
111111	1	1	0	1	1	T1 = E1
111212	6	2	0	0	0	T2 = E2
112112	3	1	0	1	0	T3 = E3
112233	2	1	0	0	1	T4 = E4
112332	3	1	1	1	0	T5 = E5
121323	3	0	2	0	0	T6 = E6
123123	1	0	1	0	0	T7 = E7
111122	6	4	0	2	3	T8 = 2 E1 + E8
112323	6	1	2	0	0	T9 = $\frac{1}{2}$ E2 + E6
	4	2	0	0	1	E8
	4	2	1	1	1	E9
	9	3	2	1	3	E10



1/8 všetkých farbení

Výsledky:

- koeficienty ako nástroj na štúdium farbení
- indikátory (p, r) -súvislosti a (p, r) -2-súvislosti
- návrh a implementácia algoritmu na výpočet koeficientov
- výpočet koeficientov pre planárne kubické k -póly do ~ 30 vr.
- analýza teoretických a empirických obmedzení pre koeficienty
- hypotéza o minimálnom pomere počtu farbení v 5-póloch

Ďalší výskum:

- dôkaz hypotézy?
- systematickejšia analýza pre obmedzenia dané dokresleniami na planárny snark